

# A method for spam blocking

Les Hatton  
CISM, University of Kingston\*

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## Abstract

Spam (unsolicited e-mail) can be split into two forms. The first is where the address of the spam is explicit and accurate, in other words the originator of the e-mail is prepared to be visible and contactable directly. This is very similar to the unsolicited mail which arrives by normal mail delivery methods. The second and by far the most common, is where the originator masquerades as a legitimate company in order to sell products or services through a third address embedded in the body of the e-mail either as an e-mail address or as a web reference.

This paper proposes a method of adjusting Mail sending and receiving software to prevent the second most numerous class of spam whereby a legitimate user's e-mail identity is hijacked.

## 1 Overview

This paper splits spam into two forms:-

1. Spam in which the originating address is genuine and belongs to the person proposing a product or service.
2. Spam in which the originating address is genuine but does not belong to the person proposing a product or a service. Instead the address has been hijacked to attempt to get through spam blocking procedures and the real response address is hidden within the body of the message either as an e-mail or as a web-site reference. The vast majority of spam received by the author (more than 99%) falls into this category and usually involves selling products which would be considered dubious at best and very likely offensive to most.

This paper deals specifically with a method to block the second category above by an encyphered verification system built on top of the normal e-mail protocol. Blocking the first category will not be considered further here.

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\*L.Hatton@kingston.ac.uk, lesh@leshatton.org

## 2 The method

Let  $S$  be the genuine e-mail address of a sender and  $R$  be the genuine e-mail address of a receiver. Typically a spammer will masquerade as  $S$  in representing itself to  $R$  in order to introduce a third contact point  $T$ , the address of the service or product being foisted on the receiver. The masquerading is done to make it more difficult to detect using simple blacklists. Furthermore, let  $K(P,M1,M2)$  be the one-way encyphered address of two e-mail addresses  $M1$ ,  $M2$  and a password or phrase,  $P$ . In this sense one-way means that the mapping  $P,M1,M2 \rightarrow K(P,M1,M2)$  for a given mail-address or addresses is as is used for example in verifying passwords on operating systems such as Linux.

Let  $P_s$ ,  $P_r$  be the sender's and receiver's keys. The method works by generating an encyphered key using the From: and To: e-mail addresses and some password  $P$  locally defined and known only to the sender as follows.

### 2.1 Typical transaction

**e-mail sender** Prefix  $K(P_s,S,R)$  to the address in From: field.

**e-mail receiver**

```
if key is absent in From: field
    Process as normal e-mail. Ultimately, reject.
else
    Generate  $K(P_r,S,R)$  and check against each key
    if 1 key in From:
        if match
            Read mail
        else
            Add  $K(P_r,S,R)$  to existing key in From: field and return to sender
        endif
    elseif 2 keys in From:
        if match
            remove matching key and return to sender
        else
            trash mail
        endif
    else
        trash mail
    endif
endif
```

Here, matching a key means encoding the passphrase and the  $S$  and  $R$  e-mail addresses and generating the same key.

An example is shown in Figure 1. A simple transaction requires 3 mails rather than one but this would be a small overhead if it prevented category 2 spam as defined above.

Figure 1: An example of a simple sender - receiver mail transaction.

Note also that this method is immune to man-in-the-middle attacks, so although a spammer T could extract an authentication key from R, this key would not work in any transaction between S and R. Note finally that this method works with CC and BCC files as it composes keys on the From: line only.

## **2.2 Key generation**

The quality of the key function  $K()$  need be of much lower cryptographic grade than those used to hide content. All that is necessary here is to make it too expensive to crack the keys.